

A Proof of that Linearizability is a Compositional Condition

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Compositionality

- A correctness condition CC is compositional if:
 - History H satisfies CC iff every register subhistory $H|x$ satisfies CC
- Linearizability is compositional:
 - $LIN(H) \Leftrightarrow \forall x: LIN(H|x)$
- We will prove that this is the case

Definition of Linearizability

- Remember the definition of linearizability:
 - LIN(H) iff there exists a sequential history S satisfying the following requirements:
 - (1) S is legal
 - (2) S and H are equivalent
 - (3) If $o_1 <_H o_2$ then $o_1 <_S o_2$
 - $o_1 <_H o_2$ denotes that $\text{res}(o_1) \rightarrow_H \text{inv}(o_2)$
 - $e_1 \rightarrow_H e_2$ denotes that e_1 precedes e_2 in H

The Proof

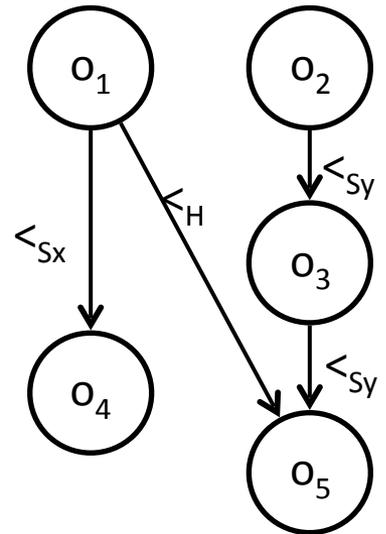
- We must show that:
 - $\text{LIN}(H) \Rightarrow \forall x: \text{LIN}(H|x)$, and
 - $\text{LIN}(H) \Leftarrow \forall x: \text{LIN}(H|x)$
- The \Rightarrow direction follows trivially
 - Exercise: Complete the proof
 - Assume $\text{LIN}(H)$, implying that there exists a sequential history S satisfying requirements (1)-(3)
 - Show that subhistory $S|x$ satisfies the requirements of a sequential history for $\text{LIN}(H|x)$

The \Leftarrow Direction

- We must show that $\text{LIN}(H) \Leftarrow \forall x: \text{LIN}(H|x)$
- Assume that the right-hand side holds:
 - For each x , there exists a sequential history S_x that satisfies the requirements of $\text{LIN}(H|x)$
 - (1) S_x is legal
 - (2) S_x and $H|x$ are equivalent
 - (3) If $o_1 <_{H|x} o_2$ then $o_1 <_{S_x} o_2$
- We will construct a sequential history S that satisfies the requirements of $\text{LIN}(H)$

Constructing Operation Graph

- Create a graph, whose vertices are operations in H , and edges are added as follows:
 - Add an edge from o to o' if $o <_{S_X} o'$
 - Add an edge from o to o' if $o <_H o'$
 - We refer to an edge as a $<_{S_X}$ edge or a $<_H$ edge (there can be zero, one, or two edges from o to o')



Constructing Sequential History S

- If the constructed graph is acyclic, then a topological sort can be performed
 - Creates a total order on operations, compatible with the partial ordering in the graph
 - History S is created directly from this total ordering
- Sequential history S created in this way satisfies the requirements for LIN(H) by construction:
 - S is legal since the total ordering in $<_{S_X}$ is legal,
 - S and H are equivalent,
 - $o_1 <_H o_2$ implies $o_1 <_S o_2$

Acyclic

- We need to show that any graph constructed as described is acyclic
 - So that the graph can be topologically sorted
 - Proof by contradiction: assume a minimal cycle exists of a certain length, and reach contradiction

Cycle of Length $n=2$

- Cycles with two operations:
 - $o_1 <_H o_2 <_H o_1$
 - Not possible, as $<_H$ is a partial order
 - $o_1 <_{Sx} o_2 <_{Sx} o_1$
 - Not possible, as $<_{Sx}$ is a total order
 - $o_1 <_H o_2 <_{Sx} o_1$
 - As o_1 and o_2 are both ops on x , then $o_1 <_H o_2$ implies that $o_1 <_{Sx} o_2$, which is contradicted in previous case

Cycle of Length $n=3$

- Cycles with three operations:
 - $O_1 <_H O_2 <_H O_3 <_H O_1$
 - Not minimal as $<_H$ is partial order
 - Similar contradiction if $<_{Sx}$ instead of $<_H$
 - $O_1 <_H O_2 <_H O_3 <_{Sx} O_1$
 - Not minimal as $<_H$ is partial order
 - In fact any cycle of length three must have two consecutive edges of same type ($<_H$ or $<_{Sx}$), and therefore cannot be minimal

Cycle of Length $n \geq 4$

- Consider a cycle of arbitrary length $n \geq 4$
- At some point in the cycle there is a section:
 - $o_1 <_H o_2 <_{Sx} o_3 <_H o_4$
 - We will show that this cycle is not minimal, as there must exist an edge $o_1 <_H o_4$
- Focus on edge $o_2 <_{Sx} o_3$, there are two cases:
 - Either $o_2 <_H o_3$, and hence $o_1 <_H o_4$ by transitivity of $<_H$
 - Or, $\text{not}(o_3 <_H o_2)$, this case is handled on next slide

Cycle of Length $n \geq 4$, case 2

- Second case, continued from previous slide:
 - $\text{not}(o_3 <_H o_2)$ implies that $\text{not}(\text{res}(o_3) \rightarrow_H \text{inv}(o_2))$
 - As \rightarrow_H is a total order, we have $\text{inv}(o_2) \rightarrow_H \text{res}(o_3)$
 - Together with $o_1 <_H o_2$, and $o_3 <_H o_4$, we have:
 - $\text{res}(o_1) \rightarrow_H \text{inv}(o_2) \rightarrow_H \text{res}(o_3) \rightarrow_H \text{inv}(o_4)$
 - Implying that $\text{res}(o_1) \rightarrow_H \text{inv}(o_4) \Rightarrow o_1 <_H o_4$
 - Hence, the cycle containing $o_1 <_H o_2 <_{S_X} o_3 <_H o_4$ is not minimal

Summary

- Cycles of all lengths have been contradicted and the graph is therefore acyclic
- It can be topologically sorted into a sequential history S that meets requirements of $LIN(H)$
- We have proven that:
 - Linearizability is compositional
 - $LIN(H) \Leftrightarrow \forall x: LIN(H|x)$